

# Text Indexing

## Lecture 00: Course Overview

Florian Kurpicz

The slides are licensed under a Creative Commons Attribution-ShareAlike 4.0 International License © ⓘ ⓘ: [www.creativecommons.org/licenses/by-sa/4.0](http://www.creativecommons.org/licenses/by-sa/4.0) | commit 6447110 compiled at 2021-10-18-09:19

# Organizational Matters

## Lectures

- Monday 10:00–11:30 (50.34, -119)
- lecture only

# Organizational Matters

## Lectures

- Monday 10:00–11:30 (50.34, -119)
- lecture only

## Project (mandatory)

- topics will be handed out 08.11.2021
- coding project and small presentation
- 20 % of the final grade

# Organizational Matters

## Lectures

- Monday 10:00–11:30 (50.34, -119)
- lecture only

## Project (mandatory)

- topics will be handed out 08.11.2021
- coding project and small presentation
- 20 % of the final grade

## Oral Exam

- 20 minutes
- 80 % of the final grade
- pizza marks content not relevant for exam



# Organizational Matters

## Lectures

- Monday 10:00–11:30 (50.34, -119)
- lecture only

## Project (mandatory)

- topics will be handed out 08.11.2021
- coding project and small presentation
- 20 % of the final grade

## Oral Exam

- 20 minutes
- 80 % of the final grade
- pizza marks content not relevant for exam



## Office Hours (Room 210)

- Monday 13:45–14:45 (lecture period)
- by appointment (otherwise)

# Materials

## Slides

- published after the lecture  
(<https://algo2.iti.kit.edu/4198.php>)

## Videos

- will be published (with  $\geq 1$  week delay)

# Materials

## Slides

- published after the lecture  
(<https://algo2.iti.kit.edu/4198.php>)

## Videos

- will be published (with  $\geq 1$  week delay)

## Additional Material

- references to literature included
- books
  - Gonzalo Navarro. *Compact Data Structures - A Practical Approach*. Cambridge University Press, 2016. ISBN: 978-1-10-715238-0
  - Enno Ohlebusch. *Bioinformatics Algorithms: Sequence Analysis, Genome Rearrangements, and Phylogenetic Reconstruction*. Oldenbusch Verlag, 2013. ISBN: 978-3000413162
- most likely no script

# Content

## Fundamentals

- tries
- suffix tree
- suffix array
- longest common prefix array
- Burrows-Wheeler transform (BWT)
- wavelet tree (+ bit vector rank/select)
- FM-index

## Compressed Indices

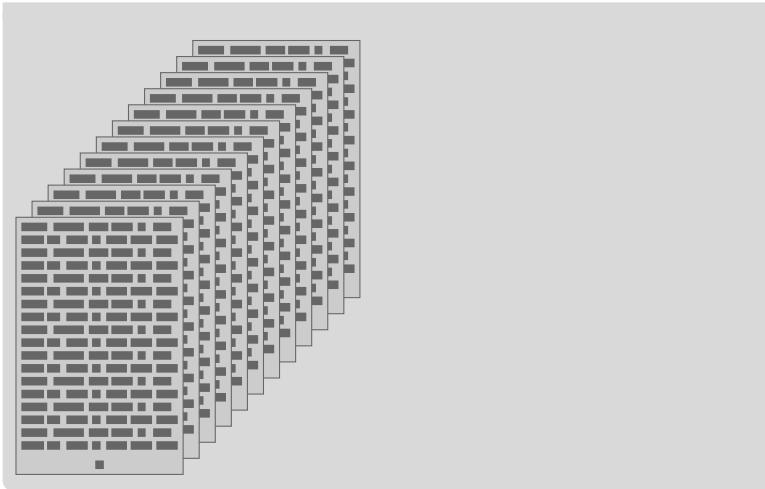
- compressing the BWT and wavelet trees
- Lempel-Ziv 77/78 compression
- LZ compression vs. BWT compression
- compressed suffix trees and suffix arrays
- r-index

## Additional Topics

- parallel construction
- different query types

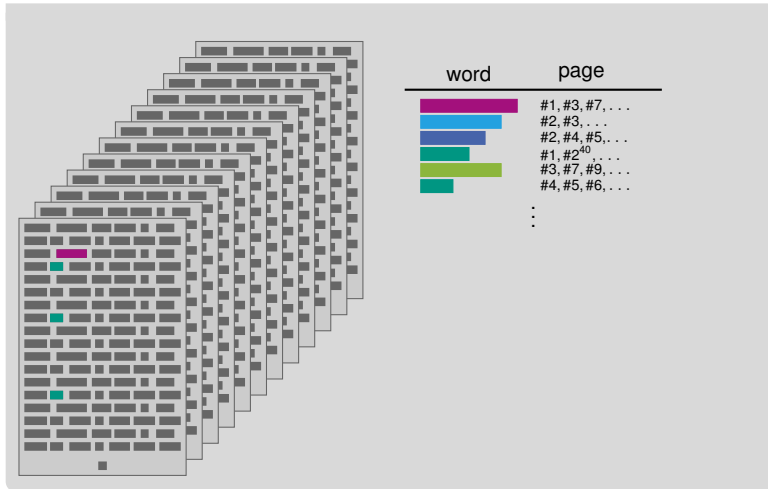


# Motivation for Text Indices



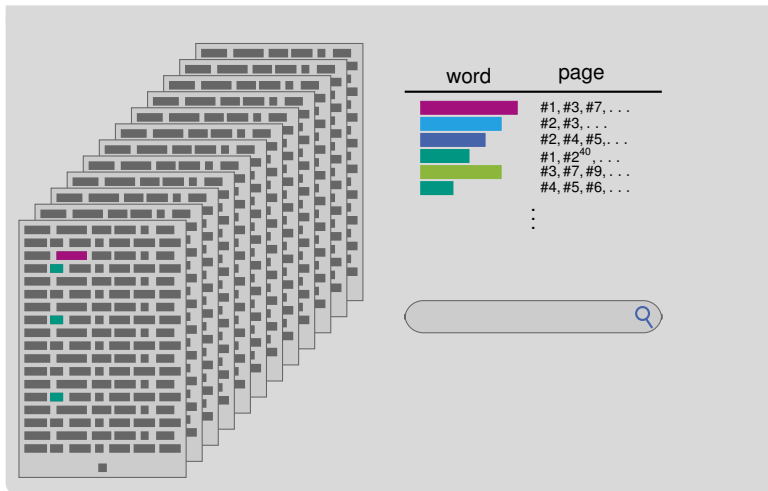
- collection of text
- scanning not feasible

# Motivation for Text Indices



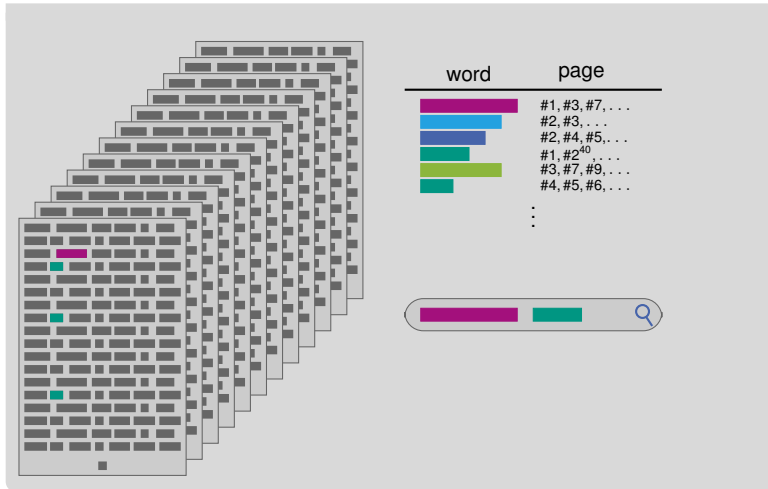
- collection of text
- scanning not feasible
- inverted index (word based)

# Motivation for Text Indices



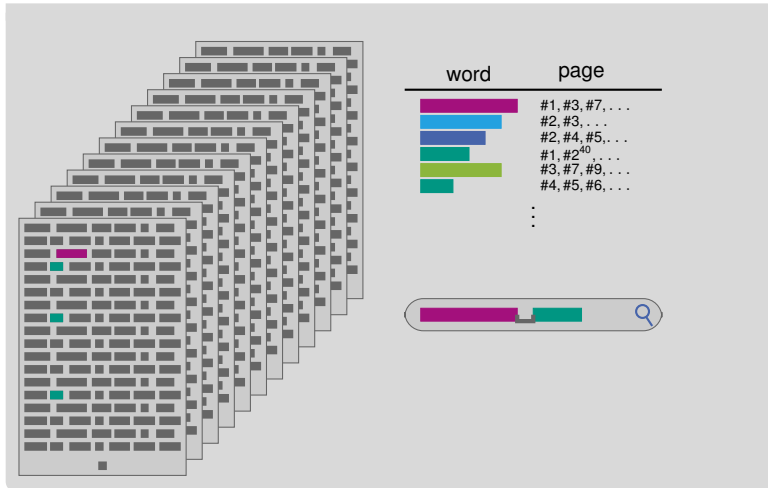
- collection of text
- scanning not feasible
- inverted index (word based)

# Motivation for Text Indices



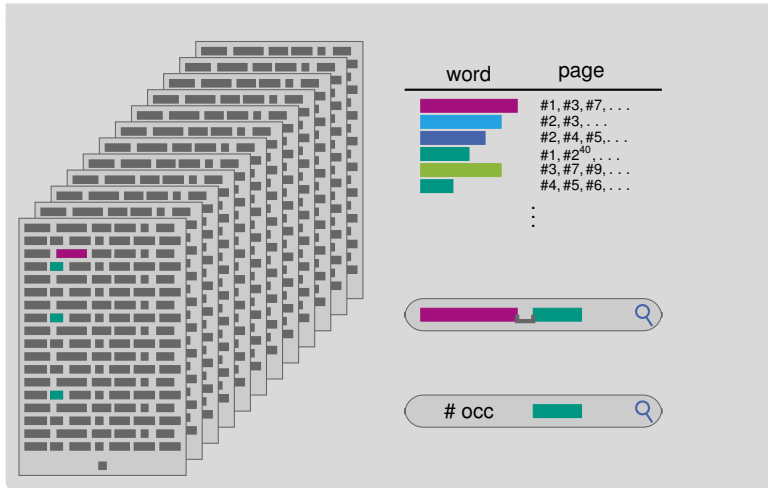
- collection of text
- scanning not feasible
- inverted index (word based)

# Motivation for Text Indices



- collection of text
- scanning not feasible
- inverted index (word based)
- phrase search

# Motivation for Text Indices









- collection of text
- scanning not feasible
- inverted index (word based)
- phrase search
- counting queries

# Motivation for Text Indices

```

GAATGCCAGTCAGCATTAAAGGCCAGGC
GGAGAGCTCAGGGCAGGTCACGTGGGA
AACTCGCATAGTGAGGGTTATCGCTCG
ACATGTTTCGTTGGGCTCTTCACTTCTT
CCGACACGAACCTCAGTTAGTTTGTTA
CCTACATCCTACCAGAGGTCGCACCTA
TGTGCCCCGGTGGTGAGAAGGAGAAGG
TGCGGATTTTCGTATTTGCAGATGCGGA
CTCGTCAGTACTTTCAGAATAACGAAT
CATGGCCTGCACGGCAAATGGCAATG
GACGCTTATAATGGACTTCGACATTTCG
AACTCGCATAGTGAGGGTTATCGGGTT
ACATGTTTCGTTGGGCTCTTCACTTTC
CCGACACGAACCTCAGTTAGTTTAGTT
TGTGCCCCGGTGGTGAGAAGGAGAAGG
CCTACATCCTACCAGAGGTCGCAGGTC
CATGGCCTGCACGGCAAATGGCAAAT
  
```

word	page
	#1, #3, #7, ...
	#2, #3, ...
	#2, #4, #5, ...
	#1, #2 <sup>40</sup> , ...
	#3, #7, #9, ...
	#4, #5, #6, ...
	⋮



- collection of text
- scanning not feasible
- inverted index (word based)
- phrase search
- counting queries
- what if there are no “words”

# Why Texts?

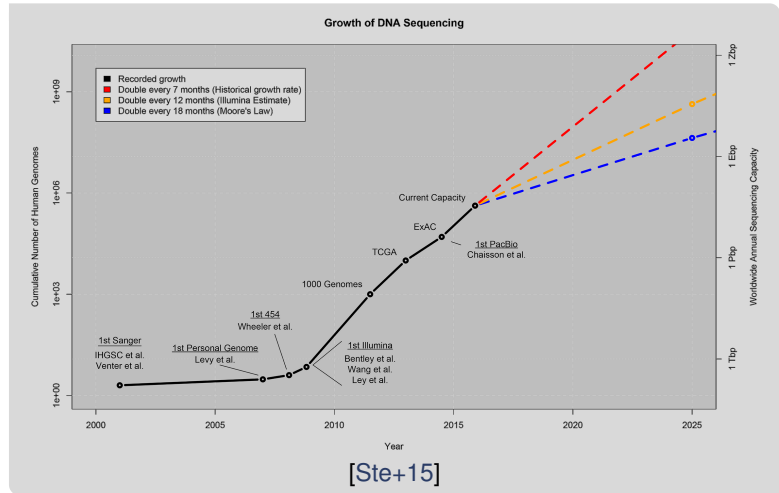
## Text is Everywhere

### Text-based Information

- Wikipedia 
- dblp 
- books 
- news articles 
- code 

### Very Important in Bioinformatics

- DNA
- proteins





# Preliminaries (1/2)

## Definition: Text

- let  $\Sigma$  be an **alphabet**
- $T \in \Sigma^*$  is a text
- $|T| = n$  is the length of the string
- $T = T[1]T[2]\dots T[n]$

# Preliminaries (1/2)

## Definition: Text

- let  $\Sigma$  be an **alphabet**
- $T \in \Sigma^*$  is a text
- $|T| = n$  is the length of the string
- $T = T[1]T[2]\dots T[n]$

## Definition: Alphabet Types

- **constant size alphabet**: finite set not depending on  $n$
- **integer alphabet**: alphabet is  $\{1, \dots, \sigma\}$  and fits into constant number of words
- **finite alphabets**: alphabet of finite size

## Preliminaries (2/2)

### Definition: Substring, Prefix, and Suffix

Given a text  $T = T[1]T[2] \dots T[n]$  of length  $n$ :

- $T[i..j] = T[i] \dots T[j]$  is called a **substring**,

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

## Preliminaries (2/2)

### Definition: Substring, Prefix, and Suffix

Given a text  $T = T[1]T[2] \dots T[n]$  of length  $n$ :

- $T[i..j] = T[i] \dots T[j]$  is called a **substring**,

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[1..i]$  is called a **prefix**, and

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

## Preliminaries (2/2)

### Definition: Substring, Prefix, and Suffix

Given a text  $T = T[1]T[2] \dots T[n]$  of length  $n$ :

- $T[i..j] = T[i] \dots T[j]$  is called a **substring**,

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[1..i]$  is called a **prefix**, and

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[i..n]$  is called a **suffix** of  $T$ .

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

## Preliminaries (2/2)

### Definition: Substring, Prefix, and Suffix

Given a text  $T = T[1]T[2] \dots T[n]$  of length  $n$ :

- $T[i..j] = T[i] \dots T[j]$  is called a **substring**,

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[1..i]$  is called a **prefix**, and

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[i..n]$  is called a **suffix** of  $T$ .

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

### Sentinel for Simplicity

Given a text  $T$  of length  $n$  over an alphabet  $\Sigma$ .

- we assume that  $T[n] = \$$  with
- $\$ \notin \Sigma$  and  $\$ < \alpha$  for all  $\alpha \in \Sigma$

## Preliminaries (2/2)

### Definition: Substring, Prefix, and Suffix

Given a text  $T = T[1]T[2] \dots T[n]$  of length  $n$ :

- $T[i..j] = T[i] \dots T[j]$  is called a **substring**,

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[1..i]$  is called a **prefix**, and

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[i..n]$  is called a **suffix** of  $T$ .

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

### Sentinel for Simplicity

Given a text  $T$  of length  $n$  over an alphabet  $\Sigma$ .

- we assume that  $T[n] = \$$  with
- $\$ \notin \Sigma$  and  $\$ < \alpha$  for all  $\alpha \in \Sigma$

# Preliminaries (2/2)

## Definition: Substring, Prefix, and Suffix

Given a text  $T = T[1]T[2] \dots T[n]$  of length  $n$ :

- $T[i..j] = T[i] \dots T[j]$  is called a **substring**,

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[1..i]$  is called a **prefix**, and

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[i..n]$  is called a **suffix** of  $T$ .

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

## Sentinel for Simplicity

Given a text  $T$  of length  $n$  over an alphabet  $\Sigma$ .

- we assume that  $T[n] = \$$  with
- $\$ \notin \Sigma$  and  $\$ < \alpha$  for all  $\alpha \in \Sigma$
- otherwise, suffix can be prefix of another suffix

1	2	3	4	5	6	7	8
a	b	b	a	a	b	b	a

- $T[1..n] = abbaabba$  and  $T[5..n] = abba$



## Preliminaries (2/2)

### Definition: Substring, Prefix, and Suffix

Given a text  $T = T[1]T[2] \dots T[n]$  of length  $n$ :

- $T[i..j] = T[i] \dots T[j]$  is called a **substring**,

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[1..i]$  is called a **prefix**, and

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

- $T[i..n]$  is called a **suffix** of  $T$ .

a	b	b	a	a	b	b	a	\$
---	---	---	---	---	---	---	---	----

### Sentinel for Simplicity

Given a text  $T$  of length  $n$  over an alphabet  $\Sigma$ .

- we assume that  $T[n] = \$$  with
- $\$ \notin \Sigma$  and  $\$ < \alpha$  for all  $\alpha \in \Sigma$
- otherwise, suffix can be prefix of another suffix

1	2	3	4	5	6	7	8
a	b	b	a	a	b	b	a

- $T[1..n] = abbaabba$  and  $T[5..n] = abba$

### Definition: Prefix-Free

A string is **prefix-free** if no suffix is a prefix of another suffix



## Bibliography

- [Nav16] Gonzalo Navarro. *Compact Data Structures - A Practical Approach*. Cambridge University Press, 2016. ISBN: 978-1-10-715238-0.
- [Ohl13] Enno Ohlebusch. *Bioinformatics Algorithms: Sequence Analysis, Genome Rearrangements, and Phylogenetic Reconstruction*. Oldenbusch Verlag, 2013. ISBN: 978-3000413162.
- [Ste+15] Zachary D Stephens., Skylar Y. Lee, Faraz Faghri, Roy H. Campbell, Chengxiang Zhai, Miles J. Efron, Ravishankar Iyer, Michael C. Schatz, Saurabh Sinha, and Gene E. Robinson. “Big Data: Astronomical or Genomical?” In: *PLOS Biology* 13.7 (July 2015), pages 1–11. DOI: [10.1371/journal.pbio.1002195](https://doi.org/10.1371/journal.pbio.1002195).